# Construction of algorithms for discrete-time quasi-birth-and-death processes through physical interpretation 

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## ACEMS

## MATHEMATICAL AND STATISTICAL FRONTIERS

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## Motivation

Algorithm construction for key matrix $\boldsymbol{\Psi}(s)$ in Stochastic Fluid Models (SFMs) includes the following three parts

- Integral expression
- Iterative scheme (central to the algorithm)
- Corresponding physical interpretation

These expressions are written in terms of the fluid generator $\mathbf{Q}(s)$.
[N. Bean, M. O'Reilly and P. Taylor. Algorithms for the Laplace-Stieltjejs transforms of return times for stochastic fluid flows. Methodology and Computing in Applied Probability, 10(3):381-408, 2008.]

## $\Psi(s)$ pictorial representation

In stochastic fluid models (SFMs) processes there is a quantity $\boldsymbol{\Psi}(s)$


## $\mathbf{G}(z)$ pictorial representation

In quasi-birth-and-death (QBD) processes there is a quantity with a similar physical interpretation


## Motivation

Here, we consider discrete-time QBDs, and

- derive summation expressions for $\mathbf{G}(s)$ with the physical interpretations similar to those of integral expressions for $\boldsymbol{\Psi}(s)$ in SFMs,
- construct corresponding iterative schemes and study them.

To do so, we use matrices $\mathbf{M}_{+}(z)$ and $\mathbf{M}_{-}(z)$, which are similar to
$\mathbf{Q}_{++}(s)$ and $\mathbf{Q}_{--}(s)$ respectively.

## Definition of the SFM

An SFM, denoted $\{(\varphi(t), X(t)): t \geq 0\}$, is a process with

- phase variable $\varphi(t)$ driven by the underlying CTMC $\{\varphi(t): t \geq 0\}$ with some finite state space $\mathcal{S}$ and generator $\mathbf{T}$,
- a level variable $X(t) \geq 0$ such that,
- when $X(t)>0$,

$$
d X(t) / d t=c_{\varphi(t)}
$$

- and when $X(t)=0$,

$$
d X(t) / d t=c_{\varphi(t)} \cdot 1\left\{c_{\varphi(t)}>0\right\}
$$

Note: $\mathcal{S}$ is partitioned as follows:
$\mathcal{S}_{+}=\left\{i \in \mathcal{S}: c_{i}>0\right\}, \quad \mathcal{S}_{-}=\left\{i \in \mathcal{S}: c_{i}<0\right\}, \quad \mathcal{S}_{0}=\left\{i \in \mathcal{S}: c_{i}=0\right\}$.

## Fluid Generator Q(s) (SFM)

$$
\mathbf{Q}(s)=\left[\begin{array}{ll}
\mathbf{Q}_{++}(s) & \mathbf{Q}_{+-}(s) \\
\mathbf{Q}_{-+}(s) & \mathbf{Q}_{--}(s)
\end{array}\right]
$$

where

$$
\begin{aligned}
& \mathbf{Q}_{++}(s)=\mathbf{C}_{+}^{-1}\left[\mathbf{T}_{++}-s \mathbf{I}-\mathbf{T}_{+0}\left(\mathbf{T}_{00}-s \mathbf{I}\right)^{-1} \mathbf{T}_{0+}\right] \\
& \mathbf{Q}_{--}(s)=\mathbf{C}_{-}^{-1}\left[\mathbf{T}_{--}-s \mathbf{I}-\mathbf{T}_{-0}\left(\mathbf{T}_{00}-s \mathbf{I}\right)^{-1} \mathbf{T}_{0-}\right] \\
& \mathbf{Q}_{+-}(s)=\mathbf{C}_{+}^{-1}\left[\mathbf{T}_{+-}-\mathbf{T}_{+0}\left(\mathbf{T}_{00}-s \mathbf{l}\right)^{-1} \mathbf{T}_{0-}\right] \\
& \mathbf{Q}_{-+}(s)=\mathbf{C}_{-}^{-1}\left[\mathbf{T}_{-+}-\mathbf{T}_{-0}\left(\mathbf{T}_{00}-s \mathbf{l}\right)^{-1} \mathbf{T}_{0+}\right]
\end{aligned}
$$

The expression $\left[e^{\mathbf{Q}(s) y}\right]_{i j}$ records the LSTs of the distribution of time for the process to have $y$ amount of fluid flowed into or out of the buffer and do so in phase $j$ given the process starts in phase $i$ and no fluid has flowed into or out of the buffer.
[ N. G. Bean, M. M. O'Reilly, and P. G. Taylor. Hitting probabilities and hitting times for stochastic fluid flows.

## Matrix $\boldsymbol{\Psi}(s)$

Let $\theta(x)=\inf \{t>0: X(t)=x\}$ be the first passage time to level $x$. For $i \in \mathcal{S}_{+}, j \in \mathcal{S}_{-}$, and $s \in \mathbb{C}$, where $\mathbb{R}(s) \geq 0,[\boldsymbol{\Psi}(s)]_{i j}$ is given by the conditional expectation

$$
[\boldsymbol{\Psi}(s)]_{i j}=E\left[e^{s \theta(x)} I\{\theta(x)<\infty, \varphi(\theta(x))=j\} \mid X(0)=x, \varphi(0)=i\right]
$$



The physical interpretation of $[\boldsymbol{\Psi}(s)]_{i j}$ is the LST of the time taken for the process to hit level $x$ for the first time and does so in phase $j$, given the process starts from level $x$ whilst avoiding levels below $x$.

## An algorithm for $\boldsymbol{\Psi}(s)$

Using $\mathbf{Q}(s)$ and physical interpretation of $\boldsymbol{\Psi}(s)$,

$$
\begin{equation*}
\boldsymbol{\Psi}_{n+1}(s)=\int_{y=0}^{\infty} e^{\mathbf{Q}_{++}(s) y}\left(\mathbf{Q}_{+-}(s)+\boldsymbol{\Psi}_{n}(s) \mathbf{Q}_{-+}(s) \boldsymbol{\Psi}_{n}(s)\right) e^{\mathbf{Q}_{--}(s) y} d y \tag{1}
\end{equation*}
$$



This can be also written as, $\boldsymbol{\Psi}_{0}(s)=\mathbf{0}$ and,

$$
\begin{equation*}
\mathbf{Q}_{++}(s) \boldsymbol{\Psi}_{n+1}(s)+\boldsymbol{\Psi}_{n+1}(s) \mathbf{Q}_{--}(s)=-\mathbf{Q}_{+-}-\boldsymbol{\Psi}_{n}(s) \mathbf{Q}_{-+}(s) \boldsymbol{\Psi}_{n}(s) \tag{2}
\end{equation*}
$$

[N. G. Bean, M. M. O'Reilly, and P. G. Taylor. Algorithms for return probabilities for stochastic fluid flows.
Stochastic Models, 21(1):149-184, 2005.]

## Definition of a discrete-time QBD

QBD is a discrete-time Markov chain, denoted $\left\{X_{t}, t \in \mathbb{N}\right\}$, on a two dimensional state space $\{(n, i): n \geq 0,1<i<m\}$, where $n$ denotes the level and $i$ the phase in state ( $n, i$ ).
One step transitions are restricted to jumps from state $(n, i)$ to ( $n^{\prime}, i^{\prime}$ ) where $n^{\prime}=n-1, n, n+1$ and $i^{\prime}$ is any phase.
The QBD has probability transition matrix $\mathbf{P}$ which is composed of square blocks $\mathbf{A}_{0}, \mathbf{A}_{+}, \mathbf{A}_{-}, \mathbf{B}$.

[G. Latouchee and V. Ramaswami. Introduction to matrix-analytic methods in stochastic modeling, volume 5.

## Probability matrix $\mathbf{P}$

$$
\mathbf{P}=\left[\begin{array}{ccccc}
\mathbf{B} & \mathbf{A}_{+} & \mathbf{0} & \mathbf{0} & \cdots \\
\mathbf{A}_{-} & \mathbf{A}_{0} & \mathbf{A}_{+} & \mathbf{0} & \cdots \\
\mathbf{0} & \mathbf{A}_{-} & \mathbf{A}_{0} & \mathbf{A}_{+} & \cdots \\
\mathbf{0} & \mathbf{0} & \mathbf{A}_{-} & \mathbf{A}_{0} & \cdots \\
\vdots & \vdots & \vdots & \vdots & \ddots
\end{array}\right]
$$

where matrices $\mathbf{B}, \mathbf{A}_{+}, \mathbf{A}_{-}, \mathbf{A}_{0}$ are square matrices of order $m$ such that, for all $i, j \in\{1 \leq i \leq m\}$,

$$
\begin{aligned}
{[\mathbf{B}]_{i j} } & =P\left(X_{t+1}=(0, j) \mid X_{t}=(0, i)\right), \\
{\left[\mathbf{A}_{+}\right]_{i j} } & =P\left(X_{t+1}=(n+1, j) \mid X_{t}=(n, i)\right), \\
{\left[\mathbf{A}_{-}\right]_{i j} } & =P\left(X_{t+1}=(n-1, j) \mid X_{t}=(n, i)\right), \\
{\left[\mathbf{A}_{0}\right]_{i j} } & =P\left(X_{t+1}=(n, j) \mid X_{t}=(n, i)\right)
\end{aligned}
$$

$\mathbf{M}_{+}(z)=\left(\mathbf{I}-\mathbf{A}_{0} z\right)^{-1} \mathbf{A}_{+} z$

$n-1$


$$
\mathbf{M}_{-}(z)=\left(\mathbf{I}-\mathbf{A}_{0} z\right)^{-1} \mathbf{A}_{-z}
$$

$$
n+1
$$



## Matrix G(z)

Let $\tau$ be the time taken to first reach level $(n-1)$. Then the $(i, j)$-th entry of the matrix $\mathbf{G}(z)$ is defined

$$
[\mathbf{G}(z)]_{i j}=E\left[z^{\tau} I\left\{\tau<\infty, X_{\tau}=(n-1, j)\right\} \mid X_{0}=(n, i)\right],
$$



The physical interpretation of $[\mathbf{G}(z)]_{i j}$ is the PGF of the time taken for the process to reach level $n-1$ for the first time and do so in phase $j$, given the process starts in level $n$ at phase $i$.

## An algorithm for $\mathbf{G}(z)$

The summation equation with a similar physical interpretation to equation (1) for $\boldsymbol{\Psi}(s)$ is

$$
\begin{equation*}
\left.\mathbf{G}_{n+1}(z)=\sum_{k=1}^{\infty} \mathbf{M}_{+}(z)^{k-1}\left(\mathbf{I}+\sum_{\ell=2}^{\infty}\left(\mathbf{M}_{+}(z) \mathbf{G}_{n}(z)\right)\right)^{\ell}\right) \mathbf{M}_{-}(z)^{k} \tag{3}
\end{equation*}
$$



This can be also written as, $\mathbf{G}_{0}^{L T}(z)=\mathbf{0}$,

$$
\begin{align*}
& \mathbf{G}_{n+1}^{L T}(z)-\mathbf{M}_{+}(z) \mathbf{G}_{n+1}^{L T}(z) \mathbf{M}_{-}(z) \\
& \quad=\left(\mathbf{I}-\mathbf{M}_{+}(z) \mathbf{G}_{n}^{L T}(z)\right)^{-1}-\mathbf{M}_{+}(s) \mathbf{G}_{n}^{L T}(z) \tag{4}
\end{align*}
$$

## Equivalence of (3) and (4)

## Lemma

## Equation

$$
\mathbf{X}=\mathbf{A X B}+\mathbf{C},
$$

for appropriately sized matrices $\mathbf{A}, \mathbf{B}$ and $\mathbf{C}$, has the unique solution given by

$$
\mathbf{X}=\sum_{k=0}^{\infty} \mathbf{A}^{k} \mathbf{C B}^{k}
$$

if and only if $\rho(\mathbf{A}) \rho(\mathbf{B})<1$, where $\rho(\cdot)$ represents the spectral radius of a given matrix.

## Convergence of $\mathbf{G}_{n}(z)$ to $\mathbf{G}(z)$

## Lemma

$\mathbf{G}_{n}^{L T}(z)$ converges to $\mathbf{G}(z)$ as $n \rightarrow \infty$.

Proof: (Outline)

- Show that $0 \leq \mathbf{G}_{n}^{L T}(z) \leq \mathbf{G}_{n+1}^{L T}(z) \leq \mathbf{G}(z)$.
- Show any arbitrary sample path for $\mathbf{G}(z)$ must be a sample path for $\mathbf{G}_{n}^{L T}(z)$ for some $n$.


## Algorithm

Input: $\mathbf{A}_{-}, \mathbf{A}_{0}, \mathbf{A}_{+}$
Set a real $\epsilon>0, z \in \operatorname{Re}>0$.
Set:
$\mathbf{M}_{+}(z)=\left(\mathbf{I}-\mathbf{A}_{0} z\right)^{-1} \mathbf{A}_{+} z$,
$\mathbf{M}_{-}(z)=\left(\mathbf{I}-\mathbf{A}_{0} z\right)^{-1} \mathbf{A}_{-} z$, and
$\mathbf{G}_{n}^{L T}(z)=\mathbf{0}$.
while $\left\|\mathbf{G}_{n+1}^{L T}(z)-\mathbf{G}_{n}^{L T}(z)\right\|_{\infty}>\epsilon$ do
Compute:

$$
C=\left(\left(\mathbf{I}-\mathbf{M}_{+}(z) \mathbf{G}_{n}^{L T}(z)\right)^{-1}-\mathbf{M}_{+}(z) \mathbf{G}_{n}^{L T}(z)\right) \mathbf{M}_{-}(z)
$$

Solve:

$$
X-\mathbf{M}_{+}(z) X \mathbf{M}_{-}(z)=C
$$

Set:

$$
\mathbf{G}_{n}^{L T}(z)=X
$$

end while
Output: $\mathbf{G}(z) \approx \mathbf{G}_{n}^{L T}(z)$

## Numerical example

Consider a QBD with $\mathbf{P}$ comprised of matrices

$$
\begin{aligned}
& \mathbf{A}_{+}=\left[\begin{array}{cccccc}
0.0151 & 0.3021 & 0 & 0 & 0 & 0 \\
0 & 0.0151 & 0.3021 & 0 & 0 & 0 \\
0 & 0 & 0.0151 & 0 & 0 & 0 \\
0 & 0 & 0 & 0.0151 & 0.3021 & 0 \\
0 & 0 & 0 & 0 & 0.0151 & 0 \\
0 & 0 & 0 & 0 & 0 & 0.0151
\end{array}\right], \\
& \mathbf{A}_{0}=\left[\begin{array}{cccccc}
0.6344 & 0.0302 & 0 & 0 & 0 & 0 \\
0.0302 & 0.6042 & 0.0302 & 0 & 0 & 0 \\
0 & 0.0302 & 0 & 0.0302 & 0 & 0 \\
0 & 0 & 0.0302 & 0.6042 & 0.0302 & 0 \\
0 & 0 & 0 & 0.0302 & 0 & 0.0302 \\
0 & 0 & 0 & 0 & 0.0302 & 0.0302
\end{array}\right],
\end{aligned}
$$

and

$$
\mathbf{A}_{-}=\left[\begin{array}{cccccc}
0.0181 & 0 & 0 & 0 & 0 & 0 \\
0 & 0.0181 & 0 & 0 & 0 & 0 \\
0 & 0 & 0.0181 & 0.9063 & 0 & 0 \\
0 & 0 & 0 & 0.0181 & 0 & 0 \\
0 & 0 & 0 & 0 & 0.0181 & 0.9063 \\
0.9063 & 0 & 0 & 0 & 0 & 0.0181
\end{array}\right] .
$$

[N. Bean, G. Latouche, and P. Taylor. Physical interpretations for quasi-birth-and-death process algorithms.

## Numerical example

Desired precision: $\epsilon=10^{-12}$

Output:

$$
\mathbf{G}=\left[\begin{array}{llllll}
0.7831 & 0.0149 & 0.0016 & 0.1084 & 0.0015 & 0.0905 \\
0.6538 & 0.0492 & 0.0030 & 0.1889 & 0.0018 & 0.1033 \\
0.0533 & 0.0016 & 0.0183 & 0.9180 & 0.0002 & 0.0087 \\
0.7426 & 0.0015 & 0.0016 & 0.1270 & 0.0022 & 0.1252 \\
0.0650 & 0.0001 & 0.0000 & 0.0040 & 0.0182 & 0.9126 \\
0.9489 & 0.0002 & 0.0000 & 0.0017 & 0.0006 & 0.0485
\end{array}\right],
$$

LT algorithm iterations: 60
Logarithmic reduction algorithm iterations: 7

LT algorithm average time: 0.015 seconds
Logarithmic reduction algorithm average time: 0.004 seconds

## Future Work

- Apply a similar idea to construct other algorithms and study them.
- Increase the complexity of the $n$-th iteration of $\mathbf{G}_{n}(z)$ and observe the outcomes.



## Thank you for listening!



## References

[N. G. Bean, M. M. O'Reilly, and P. G. Taylor. Algorithms for return probabilities for stochastic fluid flows. Stochastic Models, 21(1):149-184, 2005.]
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Pictures by Sophie Fazackerley

